LINKED LISTS

OVERVIEW

What is a linked list?

- A collection of nodes objects that have been dynamically allocated and connected to each other
- Each node object contains several pieces of information we want to store (student ID, GPA, etc.)
- Each node object has a pointer to another node object
- The pointer to the first node is stored in a head pointer
- A linked list can grow dynamically by adding more nodes



Pro: Arrays are very fast to access

We can read/write any location using array index

Con: Arrays are fixed size

Must create a new array and copy data to grow in size

Pro: Linked lists are dynamic inn size

- We can insert or delete nodes anywhere without moving data
- Con: Linked lists are slow to access
 - Must start at head and loop over linked list to find data

Review of pointer syntax:

- We declare a pointer variable using * operator
- A pointer contains the address of another variable/object int *ptr;
- We allocate memory for a variable/object using new command
- We draw this using one box pointing to another box



Review of pointer syntax:

- To access this dynamic memory we use * operator again
- The * in front of a pointer dereferences the pointer *ptr = 42;



- We can allocate memory for an array with the new command
- The array size is given inside square brackets ptr = new int[10];



Review of pointer syntax:

 To access data in dynamic array we use index notation just like we use for a static array

ptr[1] = 17;



We can also access array using * operator by adding the index to the pointer address (ugly option)
 cout << * (ptr+1);

Review of pointer syntax:

- To return memory to the O/S we use the delete command
- We should also set pointer to NULL (pointer to nothing) delete ptr;

ptr = NULL;

- To return memory for an array we add [] before pointer
- We should also set pointer to NULL (pointer to nothing) delete [] ptr; ptr = NULL;

```
CSCE 2014 - Programming Foundations II
```

Syntax for object pointers:

To store data in a linked list we must declare a Node class

```
class Node
{
    public: 
    int ID;
    float GPA;
    Node *next;
}

We simplify linked list code
by using a class with public
data and no methods
```

Syntax for object pointers:

To store data in a linked list we must declare a Node class

```
class Node
{
    public:
        int ID;
        float GPA;
        Node *next;
}
We store data in a collection
of Node attributes
```

Syntax for object pointers:

To store data in a linked list we must declare a Node class

```
class Node
{
    public:
        int ID;
        float GPA;
        Node *next;

We use the next pointer to link
this Node to another Node in list
```

Syntax for object pointers:

- We declare a Node pointer variable using * operator
- This pointer should be initialized to NULL

Node *ptr = NULL;

ptr \leftarrow We draw NULL using ground symbol

To allocate space for a Node we use the new command

ptr = new Node();



This allocates memory large enough for the Node object

- Syntax for object pointers:
 - To access fields inside a Node use use arrow -> notation ptr->ID = 123456; ptr->GPA = 3.4; ptr->next = NULL;
 ptr → 123456
 - ptr → 123456 3.4 next
 - We start at pointer variable, follow the arrow, and go to field

Syntax for object pointers:

- If the Node fields are private, we must use get and set methods in the class to access the data fields
- We use arrow -> notation before the method name ptr->setID(654321);

```
ptr->setGPA(2.7);
```

ptr->setNext(NULL);



A linked list ADT has the following operations:

- Create Initialize linked list data structure
- Insert Insert data into linked list (head, tail, sorted)
- Print Print all data in linked list
- Search Search for specific data in linked list
- Delete Remove specific value from the linked list
- Copy Copy all data from one linked list to another
- Destroy Delete linked list data structure

LINKED LISTS

OPERATION: CREATE LIST

CREATE LIST

- To create a linked list we implement Node class with the data fields we want to store and declare a head pointer
- We initialize head to NULL to create an empty linked list

```
class Node
{
   public:
    int ID;
    float GPA;
   Node *next;
}
Node *head = NULL;

The head pointer contains
the address of the first Node
in the linked list
```

LINKED LISTS

OPERATION: INSERT HEAD

 The fastest way to insert data into a linked list is to connect a new node before the head of the current list



















temp->next = head;

- When we do multiple insert head operations the data in the linked list is stored in reverse order
 - The first value inserted will be at the end of the list
 - The last value inserted will be at the head of the list



LINKED LISTS

OPERATION: INSERT TAIL

We can also insert data at the tail of the linked list

- Walk the linked list to find the last node
- Create a new node and copy data into node
- Assign pointer to connect the new node to linked list
- We must handle the special case of an empty list













while ((temp != NULL)&&
 (temp->next != NULL))
 temp = temp->next;





```
// Walk the list to end
Node * temp = head;
while ((temp != NULL)&&
    (temp->next != NULL))
    temp = temp->next;
```







head


INSERT TAIL



INSERT TAIL



head->num = 42;

```
head->next = NULL;
```

}

INSERT TAIL



LINKED LISTS

OPERATION: PRINT LIST

```
Node *ptr = head;
while (ptr != NULL)
{
    cout << ptr->ID << endl;
    cout << ptr->GPA << endl;
    ptr = ptr->next;
}
```

```
Node *ptr = head;
while (ptr != NULL)
{
    cout << ptr->ID << endl;
    cout << ptr->GPA << endl;
    ptr = ptr->next;
}
Go to the next Node in the list
```





```
Node *ptr = head;
while (ptr != NULL)
{
    cout << ptr->ID << endl;
    cout << ptr->GPA << endl;
    ptr = ptr->next;
}
```





```
Node *ptr = head;
while (ptr != NULL)
{
    cout << ptr->ID << endl;
    cout << ptr->GPA << endl;
    ptr = ptr->next;
}
```























```
Node *ptr = head;
while (ptr != NULL)
{
    cout << ptr->ID << endl;
    cout << ptr->GPA << endl;
    ptr = ptr->next;
}
```





```
void print_list(Node *ptr)
{
    if (ptr != NULL)
        {
        cout << ptr->ID << endl;
        cout << ptr->GPA << endl;
        print_list(ptr->next);
    }
}...
print_list(head)
```

```
void print_list(Node *ptr)
{
    if (ptr != NULL)
    {
        cout << ptr->ID << endl;
        cout << ptr->GPA << endl;
        print_list(ptr->next);
    }
}...
print_list(head)
```

Box method trace of recursive print_list



A linked list can also be printed in reverse order

```
void print list (Node *ptr)
{
    if (ptr != NULL)
    {
                                               Moving the recursive call
      print list(ptr->next);
                                               here will print the linked list
      cout << ptr->ID << endl;</pre>
                                               in reverse order because
      cout << ptr->GPA << endl;</pre>
                                               we print information after
                                               returning from the recursive
    }
                                               function call
}
print list(head)
```

LINKED LISTS

OPERATION: SEARCH

```
// Walk the list to desired value
Node * temp = head;
while ((temp != NULL)&&
    (temp->num != value))
    temp = temp->next;
if (temp != NULL)
    cout << value << "was found\n";
    Print message if found</pre>
```



```
// Walk the list to desired value
Node * temp = head;
while ((temp != NULL)&&
    (temp->value != 72))
    temp = temp->next;
```



```
// Walk the list to desired value
Node * temp = head;
while ((temp != NULL)&&
    (temp->value != 72))
    temp = temp->next;
```





// Walk the list to desired value
Node * temp = head;
while ((temp != NULL)&&
 (temp->value != 72))
 temp = temp->next;



```
// Walk the list to desired value
Node * temp = head;
while ((temp != NULL)&&
    (temp->value != 101))
    temp = temp->next;
    We will loop over whole list
    and stop when temp is NULL
```

LINKED LISTS

OPERATION: SORTED INSERT

SORTED INSERT

 The goal of sorted insert is to insert nodes into the list so data is always sorted in ascending (or descending) order



SORTED INSERT

Pseudo code algorithm

- Start with two pointers previous and current to head node
- Walk list until current value is greater than insertion value
- Previous value will be is less than insertion value
- Create a new node and store the insertion value
- Connect new node to list between previous and current


Implementation of sorted insert

```
// Walk the list to insertion point
Node *curr = head;
node *prev = head;
while ((curr != NULL))&&(curr->num < value))</pre>
{
     prev = curr;
                                              Stops loop when
     curr = curr->next;
                                              curr is NULL or
}
                                              num >= value
          Moves both pointers one
          Node down the linked list with
          prev one node behind curr
```



```
// Walk the list to insertion point
Node *curr = head;
node *prev = head;
while ((curr != NULL))&&(curr->num < value))
{
    prev = curr;
    curr = curr->next;
}
```



```
// Walk the list to insertion point
Node *curr = head;
node *prev = head;
while ((curr != NULL))&&(curr->num < value)) // TRUE
{
    prev = curr;
    curr = curr->next;
}
```



```
// Walk the list to insertion point
Node *curr = head;
node *prev = head;
while ((curr != NULL))&&(curr->num < value))
{
    prev = curr;
    curr = curr->next;
}
```



```
// Walk the list to insertion point
Node *curr = head;
node *prev = head;
while ((curr != NULL))&&(curr->num < value)) // TRUE
{
    prev = curr;
    curr = curr->next;
}
```



```
// Walk the list to insertion point
Node *curr = head;
node *prev = head;
while ((curr != NULL))&&(curr->num < value))
{
    prev = curr;
    curr = curr->next;
}
```



```
Node *curr = head;
node *prev = head;
while ((curr != NULL))&&(curr->num < value)) // FALSE
{
    prev = curr;
    curr = curr->next;
}
```

Implementation of sorted insert

```
// Create new node and connect to current node
Node *ptr = new Node();
ptr->num = value;
ptr->next = curr;
We know the value in the new
```

node is less than the value in the current node, so we set the next pointer to the current node



Implementation of sorted insert







where new node will be at the head of the list

LINKED LISTS

OPERATION: DELETE NODE

How can we remove data from a linked list?

- Search the linked list to find the node to delete
- Change linked list pointers to "jump over" node
- Return memory space to the operating system
- Handle special cases (delete first node, last node, empty list)





• What will this code do?



Node *prev = NULL; while ((curr != NULL) && (curr->num != 21)) { prev = curr; curr = curr->next; }

What will this code do?



89









Now we can change pointers to "jump over" deleted node



Verify code can delete node 7



Verify code can delete node 21



Verify code can delete node 3



DELETE ALL NODES

 In the destructor function we have to delete <u>all</u> nodes and give memory back to the operating system

```
Node * ptr = head;
while (ptr != NULL)
{
    head = head->next;
    delete ptr;
    ptr = head;
}
```

LINKED LISTS

HEAD AND TAIL POINTERS

HEAD AND TAIL POINTERS

 We can insert nodes at the tail much faster if we maintain a pointer to the tail of the linked list.



```
Node *ptr = new Node();
ptr->Value = "John";
ptr->Next = NULL;
Tail->Next = ptr;
Tail = ptr;
```

Clearly much faster than walking the whole linked list to find tail

HEAD AND TAIL POINTERS

 We need to adapt insert tail code to handle the special case of inserting into empty linked list

```
// Create new node
Node *ptr = new Node();
ptr->Value = value;
ptr->Next = NULL;
// Insert new node
if (Head == NULL)
    Head = ptr;
else
    Tail->Next = ptr;
Tail = ptr;
```

HEAD AND TAIL POINTERS

 We also need to modify the delete code to make sure set tail pointer correctly when the tail node is deleted

```
// Special case if we delete tail
if (curr == tail)
   tail = prev;
// Special case if we delete last node
if (head == NULL)
   tail = NULL;
```

LINKED LISTS

DOUBLY LINKED LISTS

 In order to walk the linked list in either direction, we can add a second pointer called prev to the Node object



- A doubly linked list has nice symmetry
 - We can start at head and follow next from L to R
 - We can start at tail and follow prev from R to L



- We must modify insert and delete to use both pointers
 - In some cases, code is easier because we can use prev

• For sorted insertion, we must update four pointers



• For sorted insertion, we must update four pointers



• For sorted insertion, we must update four pointers


• For sorted insertion, we must update four pointers



```
bool List::SortedInsert(int value)
{
   // Create new node
   LNode *ptr = new LNode();
   ptr->Value = value;
   ptr->Next = NULL;
   ptr->Prev = NULL;
   // Insert into empty list
   if (Head == NULL)
   {
      Head = ptr;
      Tail = ptr;
   }
```

```
// Insert before head
else if (value <= Head->Value)
{
   ptr->Next = Head;
   Head->Prev = ptr;
   Head = ptr;
}
// Insert after tail
else if (value >= Tail->Value)
{
   ptr->Prev = Tail;
   Tail->Next = ptr;
   Tail = ptr;
}
```

```
// Insert in middle
else
{
   // Walk list to deletion point
   LNode *curr = Head;
   while ((curr != NULL) && (curr->Value < value))</pre>
      curr = curr->Next;
   // Connect node to list
   ptr->Next = curr;
                                        Here is the code connecting
   ptr->Prev = curr->Prev;
                                        nodes from the slide above
   curr->Prev->Next = ptr;
   curr->Prev = ptr;
}
return true;
```

}

- For deletion from this list, we must update two pointers
 - Pointer from previous node to the next node
 - Pointer from next node to the previous node



- For deletion from this list, we must update two pointers
 - Pointer from previous node to the next node
 - Pointer from next node to the previous node



ptr->next->prev = ptr->prev;

- For deletion from this list, we must update two pointers
 - Pointer from previous node to the next node
 - Pointer from next node to the previous node



ptr->prev->next = ptr->next;

```
bool List::Delete(int value)
{
    // Walk list to deletion point
    LNode *curr = Head;
    while ((curr != NULL) && (curr->Value != value))
        curr = curr->Next;
    // Check if value was found
    if (curr == NULL)
        return false;
```

// Connect previous node to next node

```
if (curr == Head)
```

```
Head = curr->Next;
```

else

curr->Prev->Next = curr->Next;

```
// Connect next node to previous node
```

```
if (curr == Tail)
```

Tail = curr->Prev;

else

```
curr->Next->Prev = curr->Prev;
```

```
// Delete node from list
delete curr;
return true;
```

};

LINKED LISTS

SUMMARY

SUMMARY

- Linked lists can be used for applications where the amount of data varies at run time
 - We can insert and delete data without moving other data
 - We can search for specific data items in the linked list
 - We can print all of the data in the linked list
 - We can create sorted or unsorted linked lists
 - We can use tail pointers to speed up insert tail
 - We can use doubly linked lists for symmetry
- All linked list operations require that we "walk the list"
 - We can not directly access specific locations in linked list
 - We can not implement binary search on linked list